

Practical long-distance quantum communication via concatenated entanglement swapping

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&
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AK Tittel Sanders *PRA* **88** 022336 (2013)

AK Sanders *PRA* **90** 032304 (2014)

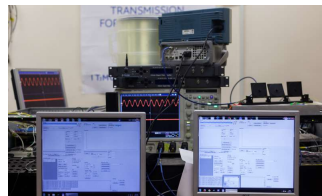
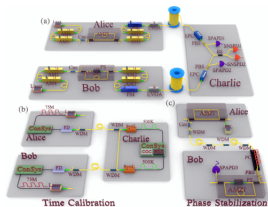
AK Sanders *JOSAB* **32** 2382 (2015)

Long-distance quantum communication



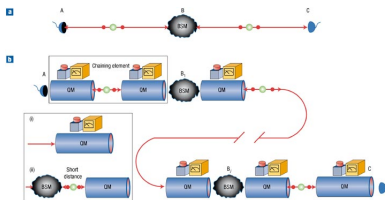
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Long-distance quantum communication



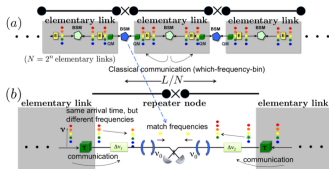
<http://goo.gl/IL95bJ>

200 km QKD: Tang et al. *PRL* (2014) 250 km QKD: Gleim et al. *OE* (2016)



Gisin & Thew *Nature Photon.* (2007)

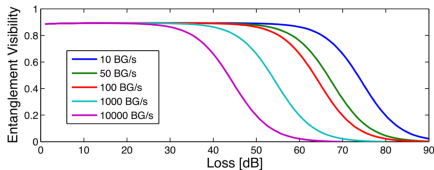
Modeling Repeaters



Errors	Approaches	Examples	Schematics	1G	2G	3G
Loss Error	Heralded Entanglement Generation (HEG)	Alice Bob Control Control		✓	✓	
	Quantum Error Correction (QEC)	$ v\rangle$ (σ) (σ) (σ) U QEC U^{-1} $ v\rangle$				✓
Operation Error	Heralded Entanglement Purification (HEP)	Alice Bob Pair 1 Pair 2 Control Control		✓		
	Quantum Error Correction (QEC)	$ v\rangle$ (σ) (σ) U QEC U^{-1} $ v\rangle$			✓	✓

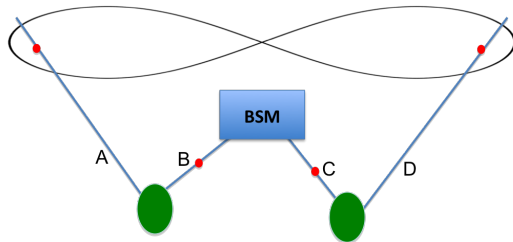
Guha et al. *PRA* (2015)

Muralidharan et al. *Sci. Rep.* (2015)



Bourgoin et al. *NJP* (2013)

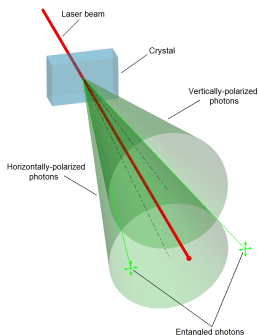
Entanglement swapping



$$|\psi^+\rangle_{AB} |\psi^+\rangle_{CD} = \frac{1}{2} \left[|\psi^+\rangle_{AD} |\psi^+\rangle_{BC} + |\psi^-\rangle_{AD} |\psi^-\rangle_{BC} + |\phi^+\rangle_{AD} |\phi^+\rangle_{BC} + |\phi^-\rangle_{AD} |\phi^-\rangle_{BC} \right]$$

Zukowski Zeilinger Horne Ekert *PRL* (1993)

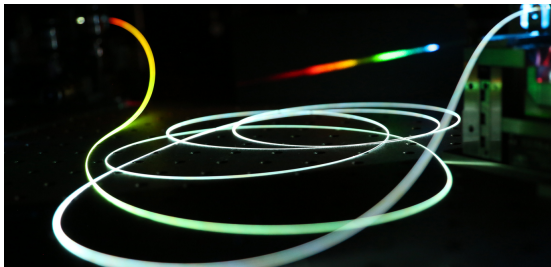
Source



<http://goo.gl/AvlyW8>

$$|\chi\rangle = e^{i\chi(\hat{a}_H^\dagger \hat{b}_H^\dagger + \hat{a}_V^\dagger \hat{b}_V^\dagger + \text{hc})} |\text{vac}\rangle = \text{sech}^2 \chi e^{i \tanh \chi (\hat{a}_H^\dagger \hat{b}_H^\dagger + \hat{a}_V^\dagger \hat{b}_V^\dagger)} |\text{vac}\rangle$$

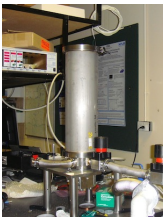
Channel



<http://goo.gl/40gB72>

$$\eta_t = e^{-(\alpha\ell + \alpha_0)/10}$$

Detector



$$\hat{\rho}_T = (1 - \exp(-\hbar\omega/KT)) \sum_{n=0}^{\infty} \exp(-n\hbar\omega/KT) |n\rangle\langle n|$$

$$\hat{\Pi}_n = |n\rangle\langle n|$$

$$\wp = \frac{(1 + \eta) \exp(-\hbar\omega/KT)}{1 - \eta \exp(-\hbar\omega/KT)}$$

$$\eta = \eta_0 \eta_t < 1$$

<http://goo.gl/EJY6wj>

$$P(q = 0|i) = (1 - \wp) [1 - \eta(1 - \wp)]^i = 1 - P(q = 1|i)$$

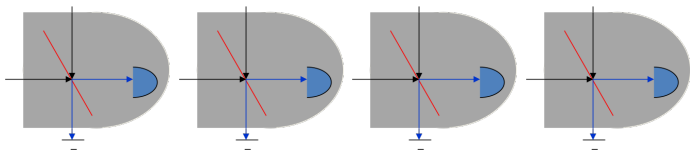
Rohde Ralph *J. Mod. Opt.* (2006)

PNR Detectors

$$\begin{aligned}
 P(q|i) &= \frac{(1-\eta)(1-\wp)}{1-\eta(1-\wp)} \left(\frac{\eta}{1-\eta} \right)^q (1-\eta)^i G(i, q; \eta, \wp) \quad \text{for } i \geq q \\
 &= \frac{(1-\eta)(1-\wp)}{1-\eta(1-\wp)} \left[\frac{1-\eta}{\eta} b(\eta, \wp) \right]^{q-i} \eta^i G(q, i; \eta, \wp) \quad \text{for } q \geq i
 \end{aligned}$$

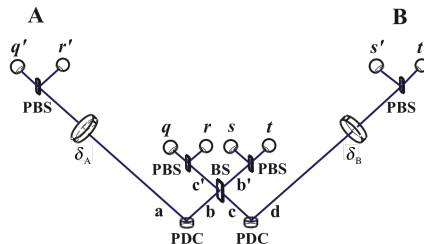
$$\begin{aligned}
 b(\eta, \zeta_{dc}) &:= \left[1 + \frac{1-\eta}{\eta\wp} \right]^{-1} \\
 G(\kappa, \lambda; \eta, \wp) &= \sum_{n=0}^{\infty} \binom{\kappa}{\lambda} \binom{\kappa - \lambda + n}{\kappa - \lambda} [b(\eta, \wp)]^n \\
 &\quad \times \left[{}_2F_1 \left(-n, -\lambda; \kappa - \lambda + 1; \frac{\eta-1}{\eta} \right) \right]^2 \quad \text{for } \kappa \geq \lambda \\
 G(\kappa, \lambda; \eta, \wp) &:= 0 \quad \text{for } \kappa < \lambda
 \end{aligned}$$

Four detectors



$$P(qrst|ijkl) = P(q|i)P(r|j)P(s|k)P(t|l)$$

Ideal single swap for 2 photons @ BSM

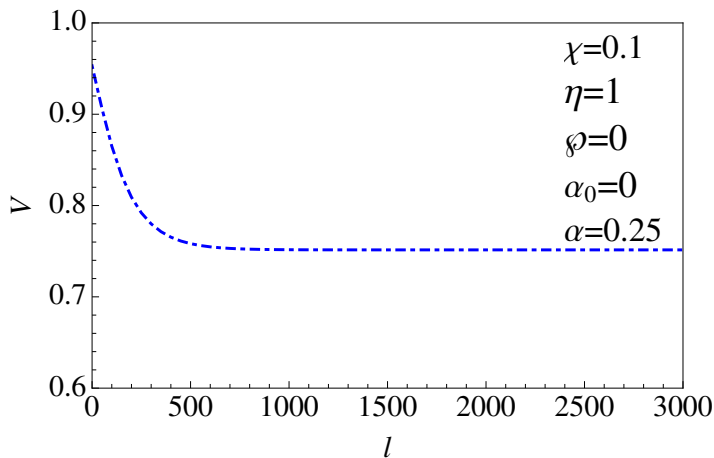


Bell state	$(qrst)$
$ \psi^+\rangle$	$(1010) \vee (0101)$
$ \psi^-\rangle$	$(0110) \vee (1001)$
$ \phi^\pm\rangle$	$(2000) \vee (0200) \vee (0020) \vee (0002)$

Conditional coincidence probability and visibility

$$Q(q'r's't'|qrst; \chi, \wp, \eta)$$
$$Q_{\text{ext}}(qrst) = \text{ext}_{q'r's't'} Q(q'r's't'|qrst; \chi, \wp, \eta)$$
$$V(\chi, \wp, \eta) = \frac{Q_{\text{max}} - Q_{\text{min}}}{Q_{\text{max}} + Q_{\text{min}}}$$

Ideal detectors



Single swap

$$|\chi\rangle_{AB}|\chi\rangle_{CD} \xrightarrow{B^{\text{inn}}} |\Xi\rangle \xrightarrow[\text{Proj}]{\text{Fock}} \frac{\Pi_{ijkl}^{\text{inn}} |\Xi\rangle}{\sqrt{P(ijkl)}} =: |\tilde{\Xi}\rangle_{ijkl}^{\text{out}}$$

$$P(ijkl) = \langle \Xi | \Pi_{ijkl}^{\text{inn}} | \Xi \rangle$$

$$|\tilde{\Xi}\rangle_{ijkl}^{\text{out}} \langle \tilde{\Xi}| \xrightarrow[\text{noisy detection}]{\text{Inner}} \rho_{qrst}^{\text{out}} = \sum P(ijkl|qrst) |\tilde{\Xi}\rangle_{ijkl}^{\text{out}} \langle \tilde{\Xi}|$$

$$\rho_{qrst}^{\text{out}} \xrightarrow[\text{polarization rotators}]{\text{Ideal counts on outer detectors given actual on inner}}$$

$$P(i'j'k'l'|qrst) = \langle i'j'k'l' | U(\delta_A)U(\delta_B)\rho_{qrst}^{\text{out}}U^\dagger(\delta_B)U^\dagger(\delta_A) | i'j'k'l' \rangle$$

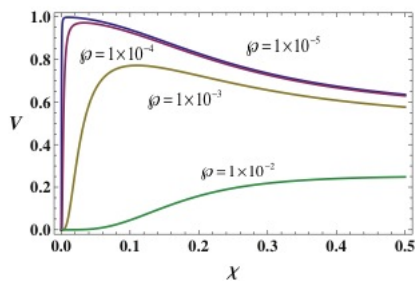
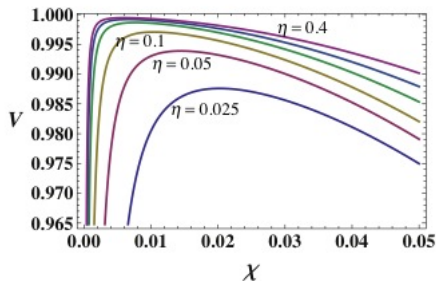
Single swap

$$\begin{aligned}
 P(ijkl|qrst) &= \frac{P(qrst|ijkl)P(ijkl)}{P(qrst)} \\
 &= P(q|i)P(r|j)P(s|k)P(t|l)P(ijkl)/P(qrst)
 \end{aligned}$$

$$Q(q'r's't'|qrst) = \sum P(q'r's't'|i'j'k'l'; \wp, \eta) P(i'j'k'l'|qrst; \chi, \wp, \eta)$$

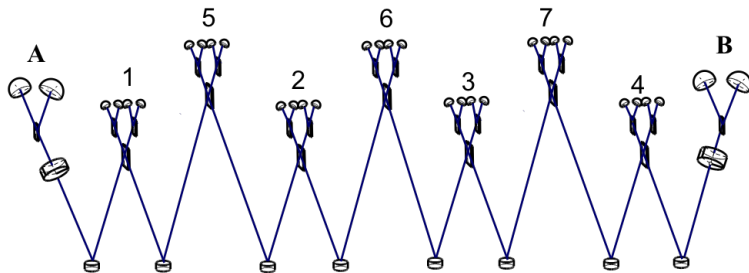
Scherer Howard Sanders Tittel *PRA* (2009)

Single-swap visibility



Scherer Howard Sanders Tittel *PRA* (2009)

N swaps with $2N - 1$ BSMs



$$Q_{\text{ext}}(qrst) = \sum_{q'r's't'}^{\text{ext}} Q(q'r's't'|qrst; \chi, \wp, \eta)$$

Khalique Tittel Sanders *PRA* (2013)

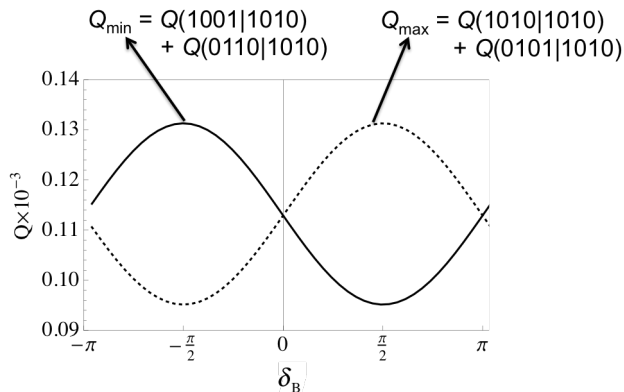
BSM connecting adjacent swaps

$$\Omega(\mu_n, \lambda_n, i_{N+n}, l_{N+n}) = \sum_{\gamma=0}^{\mu_n + \lambda_n} \binom{\mu_n + \lambda_n}{\gamma} \binom{i_{N+n} + l_{N+n} - \mu_n - \lambda_n}{i_{N+n} - \gamma} (-1)^{\mu_n + \lambda_n - \gamma}.$$

Closed-form solution

$$\begin{aligned}
P(i'j'k'l'|qrst) &= \sum_{ijkl} P(ijkl|qrst) \langle i'j'k'l' | U(\alpha) U(\delta) |\Xi\rangle_{ijkl}^{\text{out}} \langle \Xi | U^\dagger(\alpha) U^\dagger(\delta) | i'j'k'l' \rangle \\
&= \sum_{ijkl} \frac{P(qrst|ijkl)}{P(qrst)} \left(\frac{1}{\sqrt{2^{i_1+j_1+k_1+l_1} i_1! j_1! k_1! l_1!}} \frac{(\tanh \chi)^{i_1+j_1+k_1+l_1}}{\cosh^{4N} \chi} \sum_{\mu_1=0}^{i_1} \sum_{\nu_1=0}^{j_1} \sum_{\kappa_1=0}^{k_1} \sum_{\lambda_1=0}^{l_1} (-1)^{\mu_1+\nu_1} \binom{i_1}{\mu_1} \binom{j_1}{\nu_1} \right) \\
&\times \binom{k_1}{\kappa_1} \binom{l_1}{\lambda_1} \dots \frac{1}{\sqrt{2^{i_N+j_N+k_N+l_N} i_N! j_N! k_N! l_N!}} \frac{(\tanh \chi)^{i_N+j_N+k_N+l_N}}{\cosh^{4N} \chi} \\
&\times \sum_{\mu_N=0}^{i_N} \sum_{\nu_N=0}^{j_N} \sum_{\kappa_N=0}^{k_N} \sum_{\lambda_N=0}^{l_N} (-1)^{\mu_N+\nu_N} \binom{i_N}{\mu_N} \binom{j_N}{\nu_N} \binom{k_N}{\kappa_N} \binom{l_N}{\lambda_N} \Big) \\
&\times \prod_{n=1}^{N-1} \Omega(\mu_n, \lambda_n, i_{N+n}, l_{N+n}) \Omega(\nu_n, \kappa_n, j_{N+n}, k_{N+n}) \frac{\sqrt{i_{N+n}! j_{N+n}! k_{N+n}! l_{N+n}!}}{\sqrt{2^{i_{N+n}+j_{N+n}+k_{N+n}+l_{N+n}}}} \\
&\times \delta_{i_{N+n}+l_{N+n}, \mu_n+\lambda_n+i_{n+1}+l_{n+1}-\mu_{n+1}-\lambda_{n+1}} \delta_{j_{N+n}+k_{N+n}, \nu_n+\kappa_n+j_{n+1}+k_{n+1}-\nu_{n+1}-\kappa_{n+1}} \\
&\times (\nu_N + \kappa_N)! (j_1 + k_1 - \nu_1 - \kappa_1)! \sqrt{\frac{j'! k'!}{i'! l'!}} \sum_{n_a=0}^{\min[j', \nu_N + \kappa_N]} \sum_{n_d=0}^{\min[k', j_1 + k_1 - \nu_1 - \kappa_1]} \\
&\times \left(i \tan \frac{\delta_A}{2} \right)^{\nu_N + \kappa_N + j' - 2n_a} \left(\cos \frac{\delta_A}{2} \right)^{i' + j' - 2n_a} \left(i \tan \frac{\delta_B}{2} \right)^{k' + j_1 + k_1 - \nu_1 - \kappa_1 - 2n_d} \left(\cos \frac{\delta_B}{2} \right)^{l' + k' - 2n_d} \\
&\times \frac{(i' + j' - n_a)! (l' + k' - n_d)!}{n_a! n_d! (j' - n_a)! (k' - n_d)! (\nu_N + \kappa_N - n_a)! (j_1 + k_1 - \nu_1 - \kappa_1 - n_d)!} \\
&\times \delta_{i'+j', \mu_N+\nu_N+\kappa_N+\lambda_N} \delta_{k'+l', i_1+j_1+k_1+l_1-\mu_1-\nu_1-\kappa_1-\lambda_1}.
\end{aligned}$$

$$N = 3; V = 0.16 @ \delta = \pm\pi/2$$



$$\chi = 0.24$$

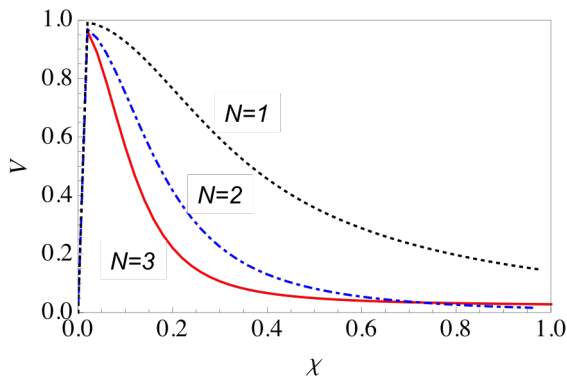
$$\eta = 0.04$$

$$\wp = 1 \times 10^{-5}$$

$$\delta_A = \pi/2$$

Khalique Sanders *PRA* (2014)

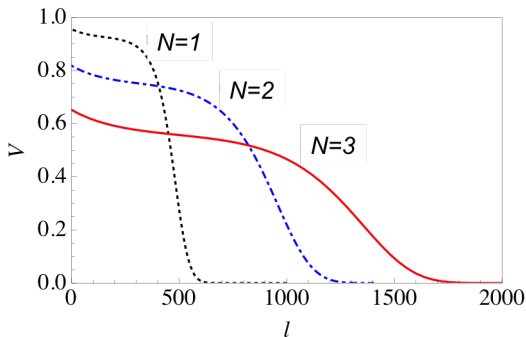
$$N \leq 3; V @ \delta_B = \pm\pi/2$$



$$\eta = 0.04$$
$$\wp = 1 \times 10^{-5}$$

Khalique Sanders *PRA* (2014)

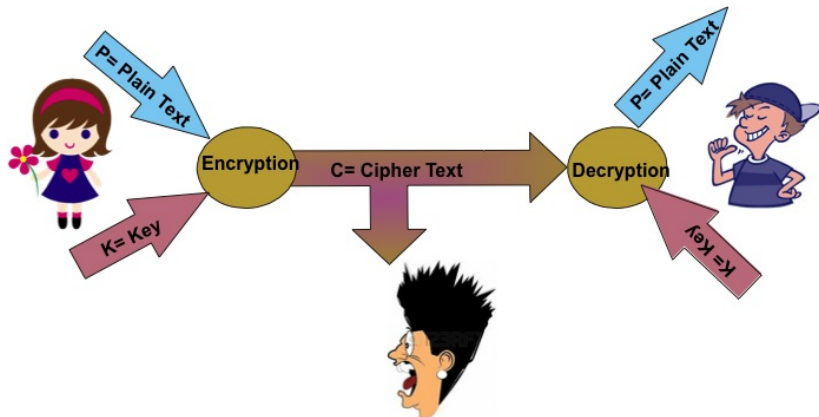
$$N \leq 3; V @ \delta_B = \pm\pi/2$$



$$\begin{aligned}\chi &= 0.1 \\ \eta_0 &= 0.70 \\ \wp &= 1 \times 10^{-5} \\ \alpha &= 0.25 \text{ dB/km} \\ \alpha_0 &= 4 \text{ dB}\end{aligned}$$

Khalique Sanders *PRA* (2014)

Quantum Cryptography



Perfect Secrecy: The Vernam Cipher

P = 0 1 1 0 1 0 0 1 0 1

K = 1 0 1 1 1 0 1 1 0 0

C = **P** $\oplus$ **K**
1 1 0 1 0 0 1 0 0 1

Impossible to break **iff** **K** is

- Long as **P**

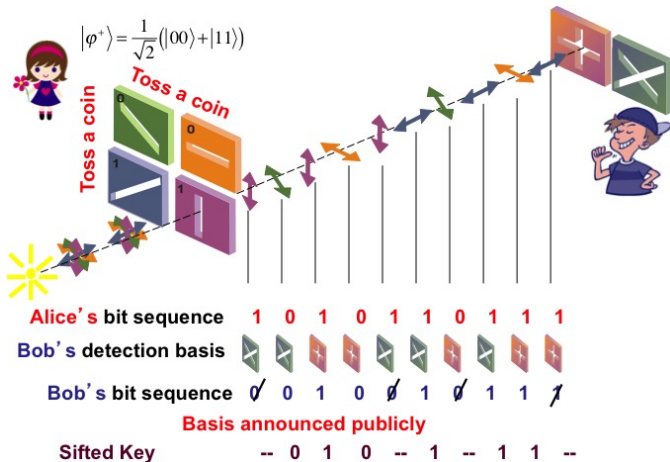


Random
Secret

- Used only once (One time pad)

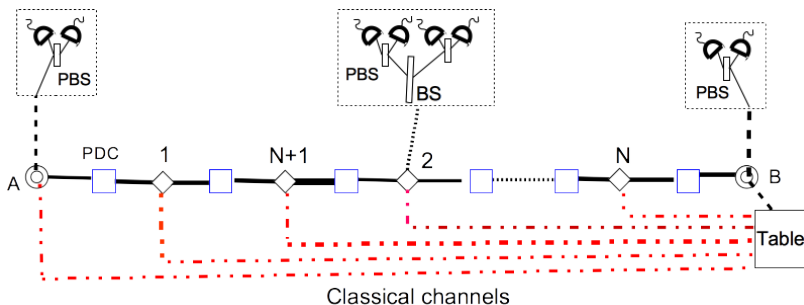
Vernam *J. Am. Inst. Elec. Eng* (1926)

Perfect Secrecy: QKD: BB84 Protocol



Bennett Brassard *IEEE* (1984)

Long-distance QKD Protocol



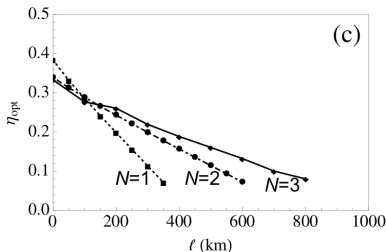
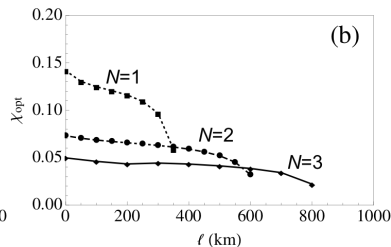
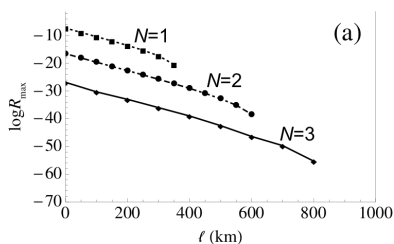
$$R = R_{\text{Shor-Preskill}} R_{\text{sifted}}$$

$$R_{\text{sifted}} = \frac{1}{2} (\chi^2)^{2N} 10^{(-\alpha l/40N)4N} (\eta^2/2)^{2N-1} \eta^2.$$

$$R_{\text{Shor-Preskill}} = 1 - \kappa H_2(Q) - H_2(Q); \quad Q = (1 - V)/2$$

Khalique Sanders *JOSA B* (2015)

Long-distance QKD



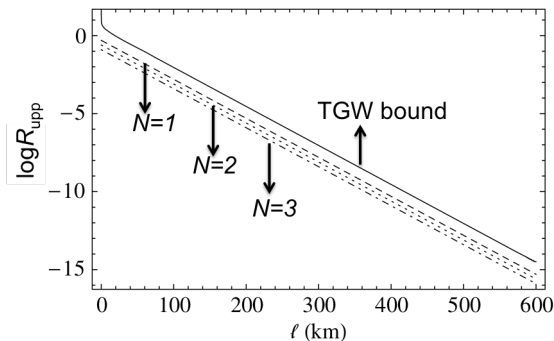
$$\wp = Ae^{B\eta}$$

For InGaAs detector

$$A = 6.1 \times 10^{-7}, B = 17$$

Khalique Sanders *JOSA B* (2015)

R_{upp} vs TGW bound for single photons



$$\wp = 0$$

$$\eta = 1$$

$$\alpha = 0.25$$

$$R_{\text{Shor-Preskill}} = 1$$

$$R = R_{\text{sifted}}$$

Khalique Sanders *JOSA B* (2015)

Truncated summation $\rightarrow$ Metropolis-Hastings sampling

With Liu Yaxiong and Zhang Pengqing (USTC)

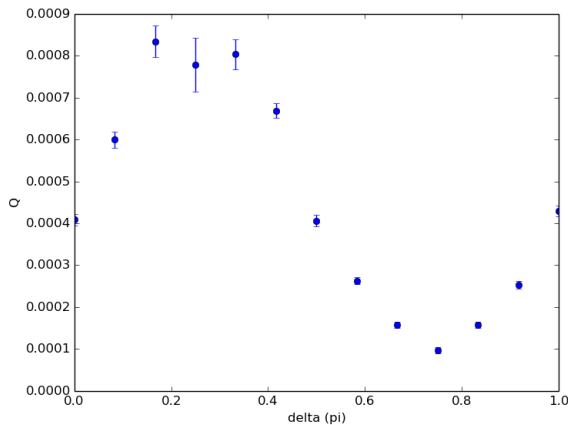
$$Q(q'r's't'|qrst) = \sum_{i'j'k'l'} P(q'r's't'|i'j'k'l') P(i'j'k'l'|qrst)$$

$$P(i'j'k'l'|qrst) = \sum_{ijkl} P(ijkl|qrst) \underbrace{\langle i'j'k'l' | U(\alpha)U(\delta) |\tilde{\Xi}\rangle_{ijkl}^{\text{out}} \langle \tilde{\Xi} | U^\dagger(\alpha)U^\dagger(\delta) | i'j'k'l' \rangle}_{A_{ijkl}^{i'j'k'l'} \text{ (hard)}}$$

- Before: $4 \times (2N - 1)$ -dimension hypercube truncation
- Now: Metropolis-Hastings sampling of $ijkl$ from distribution of $\langle A_{ijkl}^{i'j'k'l'} \rangle$ to obtain

$$Q(q'r's't'|qrst) = \langle A_{ijkl}^{i'j'k'l'} \rangle$$

Sampling result for single swap



2^{17} samples
per point
 $\chi = 0.24$
 $\eta = 0.05$
 $\varphi = 10^{-5}$
 $\delta_A = \pi/4$

Conclusions

- Developed a tractable model for practical quantum relay
- Quantified upper bounds for visibility
- In QKD, smaller N yields higher key-generation rate and larger N yields larger distances
- Next steps: Monte Carlo sampling and including quantum memory